

# Degrees of unsolvability: a two-hour tutorial

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Given a problem  $P$ , one associates to  $P$  a *degree of unsolvability*, denoted  $\deg(P)$ . Here  $\deg(P)$  is a quantity which measures the “difficulty” of  $P$ , i.e., the amount of algorithmic unsolvability which is inherent in  $P$ . We focus on two degree structures: the semilattice of Turing degrees,  $\mathcal{D}_T$ , and its completion,  $\mathcal{D}_w = \widehat{\mathcal{D}_T}$ , the lattice of Muchnik degrees. We remark that  $\mathcal{D}_w$  gives rise to a natural recursion-theoretic interpretation of Kolmogorov’s non-rigorous “calculus of problems.” We emphasize the analogy between  $\mathcal{E}_T$ , the countable sub-semilattice of  $\mathcal{D}_T$  consisting of the Turing degrees associated with recursively enumerable subsets of  $\mathbb{N}$ , and  $\mathcal{E}_w$ , the countable sublattice of  $\mathcal{D}_w$  consisting of the Muchnik degrees associated with nonempty  $\Pi_1^0$  subsets of  $\{0, 1\}^{\mathbb{N}}$ . We emphasize specific, natural degrees in  $\mathcal{D}_T$ ,  $\mathcal{D}_w$ , and  $\mathcal{E}_w$ .