

Measure-Theoretic Regularity,
Degrees of Unsolvability,
and Reverse Mathematics

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A set in Euclidean space is F_σ if it is the union of a countable sequence of closed sets. Also, the Borel sets are obtained from closed sets by repeatedly taking complements and countable unions.

Measure-theoretic regularity: Every Lebesgue measurable set includes an F_σ set of the same measure. In particular, every Borel set includes an F_σ set of the same measure.

We investigate metamathematical aspects of measure-theoretic regularity.

Degrees of unsolvability: We quantify the “descriptive complexity” or “computational strength” of the F_σ sets which are needed in order to implement measure-theoretic regularity at various levels of the effective Borel hierarchy.

Reverse mathematics: We calibrate the “logical strength” of various statements of measure-theoretic regularity.

Instead of Euclidean space, it is convenient to use the Cantor space, $\{0, 1\}^{\mathbb{N}}$.

The effective Borel hierarchy:

Let S be a subset of $\{0, 1\}^{\mathbb{N}}$. S is Σ_0^0 if and only if S is clopen. These sets are indexed in an obvious way. For each recursive ordinal α , S is $\Sigma_{\alpha+1}^0$ if and only if $S = \bigcup_{i=0}^{\infty} P_i$ where P_i is Π_{α}^0 with an index which is computable as a function of i . An index of S is an index of this computable function. S is Π_{α}^0 if its complement is Σ_{α}^0 . For limit ordinals α , S is Σ_{α}^0 if and only if S is Σ_{β}^0 for some $\beta < \alpha$.

This definition relativizes to an arbitrary Turing oracle X .

Remark. S is Borel if and only if S is $\Sigma_{\alpha}^{0,X}$ for some Turing oracle X and some $\alpha < \omega_1^X$. S is F_{σ} if and only if S is $\Sigma_2^{0,X}$ for some X .

Let $X^{(\alpha)}$ denote the α th Turing jump of X .

The *fair coin measure* on $\{0, 1\}^{\mathbb{N}}$ is given by $\mu(N_\sigma) = 1/2^{|\sigma|}$. Here σ is a *bitstring*, i.e., a finite sequence of 0's and 1's. We are writing $N_\sigma = \{Z \in \{0, 1\}^{\mathbb{N}} \mid \sigma \subset Z\}$ and $|\sigma| =$ the length of σ .

Lemma. Every $\Sigma_{\alpha+2}^{0,X}$ set includes a $\Sigma_2^{0,X^{(\alpha)}}$ set of the same measure. Conversely, every $\Sigma_2^{0,X^{(\alpha)}}$ set is $\Sigma_{\alpha+2}^{0,X}$.

Proof. For $\alpha = n$ see Steven M. Kautz, PhD thesis, Cornell, 1991. The generalization to arbitrary $\alpha < \omega_1^X$ is routine.

Definition (Nies).

1. $X \leq_{LR} Y$ if $R^Y \subseteq R^X$. Here $R^X = \{Z \mid Z \text{ is Martin-Löf random rel. to } X\}$.
2. $X \leq_{LK} Y$ if $K^Y(\tau) \leq K^X(\tau) + O(1)$ for all bitstrings τ . Here $K^X(\tau) =$ the prefix-free Kolmogorov complexity of τ relative to X .

Remark. $X \leq_T Y$ implies $X \leq_{LR} Y$,
but the converse fails badly.

Theorem. The following are equivalent.

1. $X \leq_{LR} Y$.
2. $X \leq_{LK} Y$.
3. Every $\Sigma_2^{0,X}$ set of positive measure
includes a $\Sigma_2^{0,Y}$ set of positive measure.

Theorem. The following are equivalent.

1. $X \leq_{LR} Y$ and $X \leq_T Y'$.
2. Every $\Sigma_2^{0,X}$ set includes a $\Sigma_2^{0,Y}$ set
of the same measure.

These results are due to Kjos-Hanssen
and Kjos-Hanssen/Miller/Solomon in
2005 and 2006 building on earlier work
of Dobrinen/Simpson concerning
almost everywhere domination and
measure-theoretic regularity.

Thus LR-reducibility is closely related to
measure-theoretic regularity.

Summarizing the above results, we have:

Theorem. The following are equivalent.

1. $X^{(\alpha)} \leq_{LR} Y$ and $X^{(\alpha)} \leq_T Y'$.
2. Every $\Sigma_{\alpha+2}^{0,X}$ set includes a $\Sigma_2^{0,Y}$ set of the same measure.

Recently we improved this to:

Theorem (Simpson 2009).

The following are equivalent.

1. $X^{(\alpha)} \leq_{LR} Y$ and $X \leq_T Y'$.
2. Every $\Sigma_{\alpha+2}^{0,X}$ set includes a $\Sigma_2^{0,Y}$ set of the same measure.

In particular:

Theorem (Simpson 2009).

The following are equivalent.

1. $0^{(\alpha)} \leq_{LR} Y$.
2. Every $\Sigma_{\alpha+2}^0$ set includes a $\Sigma_2^{0,Y}$ set of the same measure.

These recent improvements are based on the following technical lemmas concerning LR-reducibility.

Lemma (Simpson 2009).

Suppose $A \leq_T X$ and X is a $\Sigma_3^{0,A}$ singleton. Then $X \leq_{LR} Y$ implies $X' \leq_T A \oplus Y'$.

Cor. $X^{(\alpha)} \leq_{LR} Y$ implies $X^{(\alpha+1)} \leq_T X \oplus Y'$.

Cor. $0^{(\alpha)} \leq_{LR} Y$ implies $0^{(\alpha+1)} \leq_T Y'$.

Lemma (Simpson 2009).

If S is Σ_3^0 then

$$\begin{aligned} S^{LR} &= \text{the LR-upward closure of } S \\ &= \{Y \mid (\exists X \in S) (X \leq_{LR} Y)\} \end{aligned}$$

is again Σ_3^0 .

Cor. For each recursive ordinal α , the set

$$B_\alpha = \{Y \mid 0^{(\alpha)} \leq_{LR} Y\}$$

is Σ_3^0 .

Degrees of unsolvability:

A *mass problem* is a set of Turing oracles, $P \subseteq \{0, 1\}^{\mathbb{N}}$.

We say that P is *weakly reducible to* Q , written $P \leq_w Q$, if for each $Y \in Q$ there exists $X \in P$ such that $X \leq_T Y$.

The *weak degree of* P , written $\text{deg}_w(P)$, is the equivalence class of P under mutual weak reducibility.

The weak degrees form a lattice, \mathcal{D}_w .

Idea: P is a “problem”. The elements of P are the “solutions” of P . P is “reducible” to Q if each “solution” of Q can be used as a Turing oracle to find some “solution” of P . Two “problems” are equivalent if each is “reducible” to the other.

Remark. Muchnik introduced \mathcal{D}_w in 1963 in order to rigorously explicate Kolmogorov’s informal 1932 interpretation of intuitionism as a “calculus of problems.”

A sublattice of \mathcal{D}_W :

In recent years I have been studying a sublattice of \mathcal{D}_W which I call \mathcal{E}_W . Namely, \mathcal{E}_W is the lattice of weak degrees of mass problems associated with nonempty Π_1^0 subsets of $\{0, 1\}^{\mathbb{N}}$. It turns out that \mathcal{E}_W is extremely useful for classifying unsolvable mathematical problems. In particular, \mathcal{E}_W contains many, specific, natural degrees of unsolvability which are of interest from the viewpoint of foundations of mathematics.

Let $\mathbf{1}$ and $\mathbf{0}$ denote the top and bottom degrees in \mathcal{E}_W . A handy technical lemma is:

Lemma (Simpson 2004).

If $s = \deg_W(S)$ where S is Σ_3^0 , then $\inf(s, \mathbf{1})$ belongs to \mathcal{E}_W .

This applies to measure-theoretic regularity, because for each recursive ordinal α , the mass problem

$\{Y \mid 0^{(\alpha)} \leq_{\text{LR}} Y\} = \{Y \mid \text{every } \Sigma_{\alpha+2}^0 \text{ set includes a } \Sigma_2^{0,Y} \text{ set of the same measure}\}$ is Σ_3^0 .

Letting

$$\mathbf{b}_\alpha = \text{deg}_w(\{Y \mid 0^{(\alpha)} \leq_{\text{LR}} Y\})$$

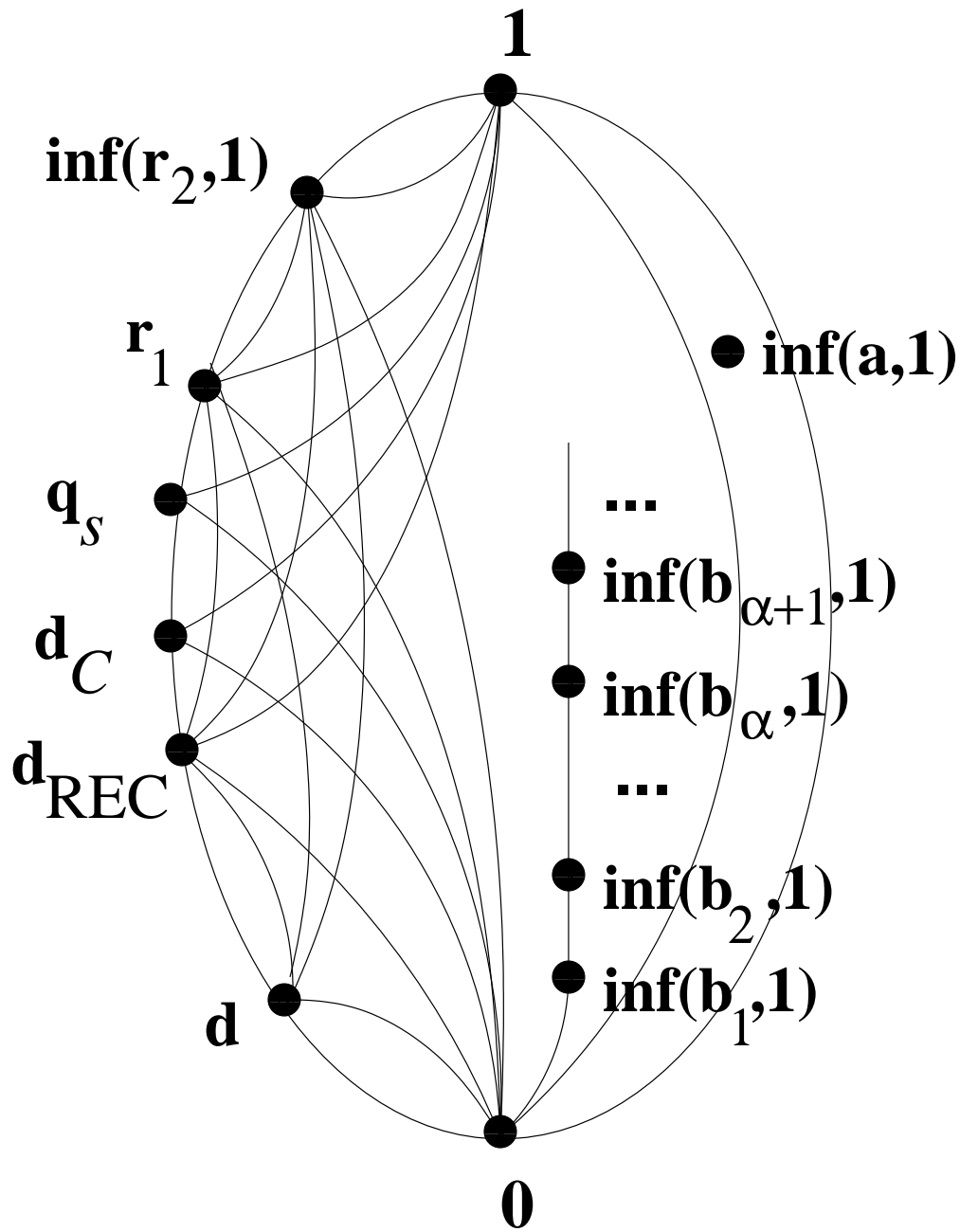
we have:

Theorem (Simpson 2009). The weak degrees $\text{inf}(\mathbf{b}_\alpha, \mathbf{1})$ belong to \mathcal{E}_w . Moreover, for $\alpha > 0$ these degrees are distinct from one another and incomparable with previously known degrees in \mathcal{E}_w including

$$\mathbf{r}_1 = \text{deg}_w(\{Z \mid Z \text{ is Martin-Löf random}\})$$

and

$$\mathbf{d} = \text{deg}_w(\{f \mid f \text{ is diagonally nonrecursive}\}).$$



A picture of \mathcal{E}_w . Here a = any r.e. degree, b = measure-theoretic regularity, r = randomness, q = Hausdorff dimension, d = diagonal nonrecursiveness.

Applications to reverse mathematics:

Applications to reverse mathematics follow, because many specific degrees in \mathcal{D}_W and \mathcal{E}_W are correlated to specific subsystems of second-order arithmetic. In particular $ACA_0 \approx \mathbf{0}'$, $WKL_0 \approx \mathbf{1}$, $WWKL_0 \approx \mathbf{r}_1$.

Recall that ACA_0 and WKL_0 are members of the “Big Five.” $WWKL_0$ is a weaker system which has been highly relevant for the reverse mathematics of measure theory.

Definition. Let M be an ω -model of RCA_0 . Let S be a set in Euclidean space.

1. S is *M-Borel* if S is $\Sigma_{\alpha}^{0,X}$ for some $X \in M$ and $\alpha < \omega_1^X$.
2. S is *M-F $_{\sigma}$* if S is $\Sigma_2^{0,Y}$ for some $Y \in M$.
3. M is an *MTR-model* if every *M-Borel* set includes an *M-F $_{\sigma}$* set of the same measure.

Note: MTR = “measure-theoretic regularity.”

Lemma. M is an MTR-model if and only if $(\forall X \in M) (\forall \alpha < \omega_1^X) (\exists Y \in M) (X^{(\alpha)} \leq_{\text{LR}} Y)$.

Cor. If $M \models \text{ATR}_0$ then M is an MTR-model.

Using our new results on LR-reducibility, we can build interesting MTR-models.

Theorem (Simpson 2009). We can find MTR-models M_1, M_2, M_3, M_4 satisfying $\text{RCA}_0 + \neg \text{WWKL}_0$ and $\text{WWKL}_0 + \neg \text{WKL}_0$ and $\text{WKL}_0 + \neg \text{ACA}_0$ and $\text{ACA}_0 + \neg \text{ATR}_0$ respectively.

Moreover, these models can be made to satisfy measure-theoretic regularity at all levels of the Borel hierarchy corresponding to countable well-orderings with a small amount of transfinite induction.

Tentative conclusion:

The reverse mathematics of measure-theoretic regularity appears to be somewhat orthogonal to the Gödel hierarchy including the “Big Five” subsystems of second-order arithmetic.

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