

Computability, Unsolvability, Randomness

Math 497A: Homework #11

Stephen G. Simpson

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For $f, g \in \mathbb{N}^{\mathbb{N}}$ say that f is *majorized* by g if $f(n) < g(n)$ for all n .

1. If $P(f, g, -)$ is a Π_1^0 predicate, prove that the predicate

$$Q(g, -) \equiv \exists f (P(f, g, -) \wedge f \text{ is majorized by } g)$$

is again Π_1^0 .

Note: This is a generalization of the Magic Lemma, Lemma 48.3 in the Lecture Notes. You can prove it by imitating the the proof of the Magic Lemma.

2. (a) Show that the result of Problem 1 holds if we replace Π_1^0 by Σ_2^0 .
(b) Show that the result does not hold if we replace Π_1^0 by Π_2^0 .
In fact, we can find a Π_2^0 predicate $P(X, -)$ with X ranging over $2^{\mathbb{N}}$ such that the predicate $\exists X P(X, -)$ is not arithmetical, i.e., it is not Π_n^0 or Σ_n^0 for any n .
3. Let $P \subseteq 2^{\mathbb{N}}$ be Π_1^0 . Let $\Phi(X, n)$ be a partial recursive functional such that $\Phi(X, n) \downarrow$ for all $X \in P$ and all n . Find a total recursive function $g(n)$ which majorizes $\Phi(X, n)$ for all $X \in P$ and all n .
4. An oracle X is said to be *hyperimmune-free* (sorry for the awkward terminology) if each $f \leq_T X$ is majorized by some recursive function.

Note: This is another example of a “lowness property” of X .

- (a) Let $P \subseteq 2^{\mathbb{N}}$ be nonempty and Π_1^0 . Prove that there exists $X \in P$ such that X is hyperimmune-free. This result is known as the Hyperimmune-Free Basis Theorem.

Hint: Use Π_1^0 approximation as in the Low Basis Theorem.

- (b) Deduce that we can find a random X which is hyperimmune-free.

5. Prove that if $0 <_T X \leq_T 0'$ then X is not hyperimmune-free.

Note: This prevents us from combining the Low Basis Theorem and the Hyperimmune-Free Basis Theorem into one theorem.

Hint for the proof: By Post's Theorem X is Δ_2^0 . Deduce that the singleton set $\{X\}$ is Π_2^0 . Use this to find $f \equiv_T X$ such that the singleton set $\{f\}$ is Π_1^0 . If such an f is majorized by a recursive function, use the result of Problem 1 to show that f is recursive.

6. (Extra Credit)

- (a) Prove that if X is 2-random then X is not hyperimmune-free.
 (b) What if we assume only that X is weakly 2-random?

7. (a) Prove that if Y is nonrecursive then $\mu(\{X \in 2^{\mathbb{N}} \mid Y \not\leq_T X\}) = 1$.
 (b) Deduce that for each nonrecursive Y we can find a random X such that $Y \not\leq_T X$.
 (c) More generally, prove the following. Given a sequence of nonrecursive oracles $Y_i, i = 0, 1, 2, \dots$, we can find an X which is n -random for all n and such that $Y_i \not\leq_T X$ for all i .

Note: It can be shown that for all Y we can find a random X such that $Y \leq_T X$. In fact, each Turing degree $\geq \mathbf{0}'$ contains a random X . However, this does not hold for weakly 2-random X 's, because all such X 's are Turing incomparable with $0'$.

8. (a) Assume that $P \subseteq 2^{\mathbb{N}}$ is Π_1^0 and

$$\neg \exists X (X \in P \wedge X \text{ is recursive}).$$

Find a nonrecursive Y such that

$$\neg \exists X (X \in P \wedge X \leq_T Y).$$

Hint: Use finite approximation.

- (b) Find a nonrecursive Y such that

$$\neg \exists X (X \text{ is random} \wedge X \leq_T Y).$$

Hint: Use the fact that $\{X \mid X \text{ is random}\}$ is the union of a sequence of Π_1^0 sets.